The Digital Abstraction

## Review

- Discretize matter by agreeing to observe the lumped matter discipline



## Lumped Circuit Abstraction

## - Analysis tool kit: KVL/KCL, node method, superposition, Thévenin, Norton (remember superposition, Thévenin, Norton apply only for linear circuits)

## Today

## Discretize value $\longrightarrow$ Digital abstraction

## Interestingly, we will see shortly that the tools learned in the previous three lectures are sufficient to analyze simple digital circuits

## Reading: Chapter 5 of Agarwal \& Lang

## But first, why digital?

## In the past

## Analog signal processing



## By superposition,

$$
V_{0}=\frac{R_{2}}{R_{1}+R_{2}} V_{1}+\frac{R_{1}}{R_{1}+R_{2}} V_{2}
$$

If $R_{1}=R_{2}$,

$$
V_{0}=\frac{V_{1}+V_{2}}{2}
$$

## The above is an "adder" circuit.

## Noise Problem


add noise on
this wire $\mathbf{~ м и м m ~}$

... noise hampers our ability to distinguish between small differences in value e.g. between 3.1V and 3.2V.

# Value Discretization 

## Restrict values to be one of two

## HIGH

5 V
TRUE
1

LOW
OV
FALSE
0
...like two digits 0 and 1

## Why is this discretization useful?

> (Remember, numbers larger than 1 can be represented using multiple binary digits and coding, much like using multiple decimal digits to represent numbers greater than 9. E.g., the binary number 101 has decimal value 5.)

## Digital System noise <br> 

## With noise



Cite as: Anent Agarwal and Jeffrey Lang, course materials for 6.002 Circuits and Electronics, Spring 2007. MIT OpenCourseWare (http://ocw.mit.edu/), Massachusetts Institute of Technology. Downloaded on [DD Month YYYY].

### 6.002 Fall 2000 Lecture 4

## Digital System

## Better noise immunity Lots of "noise margin"

## For "1": noise margin 5 V to $2.5 \mathrm{~V}=2.5 \mathrm{~V}$ For " 0 ": noise margin $0 V$ to $2.5 \mathrm{~V}=2.5 \mathrm{~V}$

## Voltage Thresholds and Logic Values



## But, but, but What about 2.5V?

Hmmm... create "no man's land" or forbidden region

## For example,



$$
\begin{aligned}
& " 1 " \rightarrow V_{H} \rightarrow 5 \mathrm{~V} \\
& { }^{" 0 "} \rightarrow \mathrm{OV} \rightarrow \mathrm{~V}_{\mathrm{L}}
\end{aligned}
$$

## But, but, but ... Where's the noise margin?

 What if the sender sent 1: $V_{H}$ ? Hold the sender to tougher standards!

## But, but, but ... Where's the noise margin?

 What if the sender sent 1: $V_{H}$ ? Hold the sender to tougher standards!
"1" noise margin: $V_{I H}-V_{O H}$
" 0 " noise margin: $V_{I L}-V_{O L}$


Digital systems follow static discipline: if inputs to the digital system meet valid input thresholds, then the system guarantees its outputs will meet valid output thresholds.

## Processing digital signals

## Recall, we have only two values -

## $1, \mathbf{0} \Rightarrow$ Map naturally to logic: $T, F$ $\Rightarrow$ Can also represent numbers

## Processing digital signals

## Boolean Logic

## $\Rightarrow$ If $X$ is true and $Y$ is true

 Then $\mathbf{Z}$ is true else $\mathbf{Z}$ is false.$$
\begin{aligned}
\Rightarrow Z & =X \text { AND } Y \longrightarrow \underbrace{Z=X \cdot Y}_{\text {Boolean equation }} \xrightarrow{ } \quad \begin{array}{c}
X, Y, Z \\
\text { are digital signals } \\
\text { "0" "1" }
\end{array}
\end{aligned}
$$



## $\Rightarrow$ Truth table representation:



Enumerate all input combinations

## Combinational gate abstraction

## - Adheres to static discipline Outputs are a function of inputs alone.

## Digital logic designers do not

 have to care about what is inside a gate.


Noise


$$
z=x \cdot y
$$

### 6.002 Fall 2000 <br> Lecture 4

## Examples for recitation



## $z=x \cdot y$

## In recitation...

## Another example of a gate

If ( $A$ is true) OR ( $B$ is true)
then $C$ is true
else $C$ is false
$\Rightarrow C=A+\underbrace{B \quad \begin{array}{c}\text { Boolean equation } \\ O R\end{array}}$


OR gate

## More gates



Inverter

$z=\overline{x \cdot y}$

## Boolean Identities

$$
\begin{aligned}
x \cdot 1 & =x \\
x \cdot 0 & =x \\
x+1 & =1 \\
x+0 & =x \\
\overline{1} & =0 \\
0 & =1 \\
A B+A C & =A \cdot(B+C)
\end{aligned}
$$

## Digital Circuits

## Implement: output $=A+\overline{B \cdot C}$



